Simplified POMDP Algorithms with Performance Guarantees

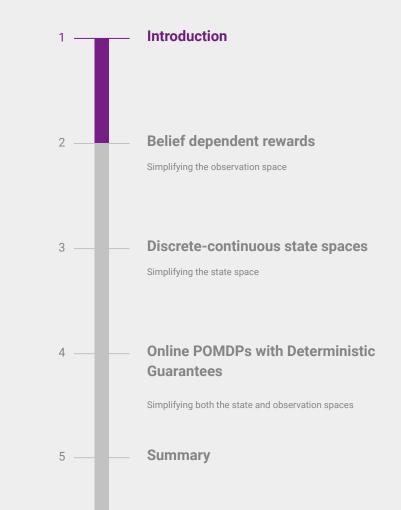
Moran Barenboim

Supervisor: Assoc. Prof. Vadim Indelman





Simplified POMDP Algorithms with Performance Guarantees





Introduction

Sequential decision-making under uncertainty

Examples include,



Indoor

Upenn

Urban IROS 2013 workshop Autonomous Cars Georgia Tech





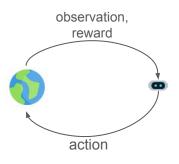
Introduction - Formalism

Partially Observable Markov Decision Process (POMDP)

$$\mathcal{M} = \langle \mathcal{X}, \mathcal{A}, \mathcal{Z}, \mathcal{P}_0, \mathcal{P}_T, \mathcal{Z}, \mathcal{R}, \mathcal{T} \rangle$$

- State $x_t \in \mathcal{X}$
- Action $a_t \in \mathcal{A}$
- Observation $z_t \in \mathcal{Z}$
- Reward $r(x_t, a_t)$
- Transition function $P_T(x_{t+1}, x_t, a_t) = \mathbb{P}(x_{t+1}|x_t, a_t)$
- Observation function $P_z(z_t, x_t) = \mathbb{P}(z_t|x_t)$
- Prior distribution (also prior belief) b_0
- \blacksquare Horizon $\mathcal T$





Introduction - Formalism

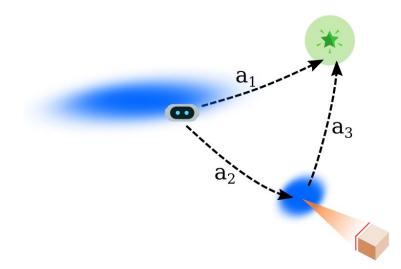
- History $H_t = \{b_0, a_0, z_1, a_1, \dots a_{t-1}, z_t\}$
- Belief $b_t = \mathbb{P}(x_t \mid H_t)$
- Policy $\pi_t(b_t)$
- Value function $V^{\pi}(b_t)$, where,

$$egin{aligned} &\mathcal{V}^{\pi_t}(b_t) = \mathbb{E}_{z_{t+1:\mathcal{T}}}\left[\sum_{i=t}^{\mathcal{T}} r(b_i,\pi_i)
ight] \ &= r(b_t,\pi_t) + \mathbb{E}_{z_{t+1}}\left[\mathcal{V}^{\pi_t}(b_{t+1})
ight] \end{aligned}$$



Introduction - Formalism

The (optimal) solution for a POMDP optimally trades off information-gathering actions versus other actions.





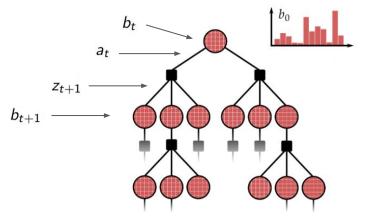


Introduction - Solutions

We will be focusing on online tree search methods

- Each node represents a belief
- Each edge represents an action or an observation
- Given a prior belief, the posterior belief is calculated via probabilistic inference

$$b_{t+1} = \psi(b_t, a_t, z_{t+1})$$





Introduction - Solutions

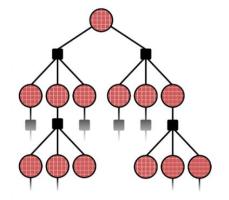
How to get an exact solution?

Given a POMDP definition, construct a tree of all states, actions and observations

Problem?

Size of the tree - $O(\mathcal{T}^{|O| \times |A|})$

Only relevant for very small POMDPs







Introduction - Solutions

Approximate planners:

	Planning efficiency	Aware of state uncertainty	Optimal (in some sense)
Gradient-based, open-loop	Yes	No	No
Deterministic approximations	Yes	No	No
Monte-Carlo Sampling	Yes	Yes	Yes

POMCP, DESPOT, POMCPOW, PFT-DPW, AdaOPS...



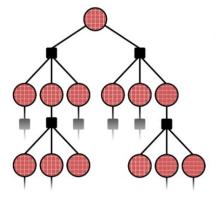


Introduction - Our Approach

In this research, we derive a framework instead of solving the original POMDP, considers a simplified version of that POMDP.

 $\mathcal{M} \longrightarrow \bar{\mathcal{M}}$

Then, we aim at deriving a mathematical relationship between the solution of the simplified, and the theoretical POMDP.



 $\bar{V}(b)$ V(b)





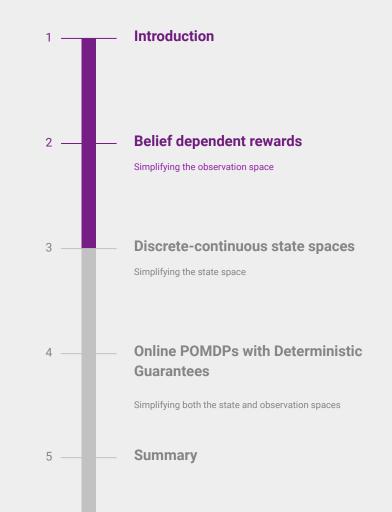
Simplified POMDP Algorithms with Performance Guarantees

Published:

Technology

M. Barenboim and V. Indelman. Adaptive information belief space planning. In the 31st International Joint Conference on Artificial Intelligence and the 25th European Conference on Artificial Intelligence (IJCAI-ECAI), July 2022

> Autonomous Navigation and Perception Lab



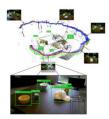
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Belief Based Rewards - Motivation

A generalization of POMDPs (limited to state-based rewards).

Supports explicit reasoning of uncertainty, e.g.,

- Pose uncertainty (of the robot, other agents, etc.)
- Map representation
- Semantic uncertainty



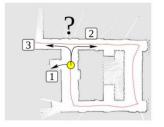


Figure: Stachniss et al.

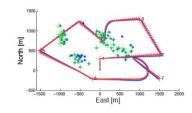


Figure: Indelman et al.





Technology

Belief Based Rewards - Motivation

Information-Theoretic functions may be used as uncertainty measures, commonly used as reward functions. E.g.,

Differential entropy

$$\mathcal{H}(x_t) = -\int_x b_t \cdot \log(b_t) dx$$

- Information Gain
- Mutual information
- Kullback-Leibler divergence
- and more...



Belief Based Rewards - Introduction

We focused on entropy as an information-theoretic reward function,

$$r(b_t, a_t, b_{t+1}) = \omega_1 \mathbb{E}_{s \sim b_{t+1}} [r(s, a_t)] + \omega_2 \mathcal{H}(b_{t+1}),$$

a weighted sum of state-dependent reward and entropy (discrete or continuous)



Belief Based Rewards - The Challenge

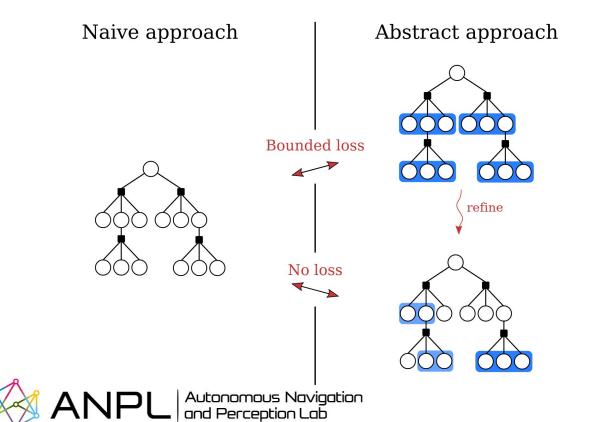
The difficulty - Information theoretic functions are generally intractable

And even approximations are computationally difficult - $O(|\mathcal{X}|^2)$ - for every reward calculation



ECHNION

of Technology



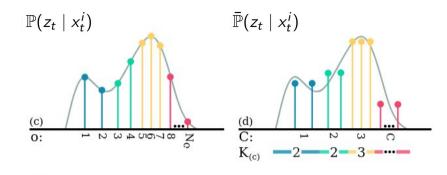


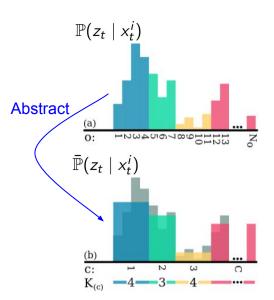
We introduced an abstract observation model,

$$\bar{\mathbb{P}}(z_t^j \mid x_t^i) = rac{\sum_k^{\kappa} \mathbb{P}(z_t^k \mid x_t^k)}{\kappa}$$

• Aggregates a set of K observations

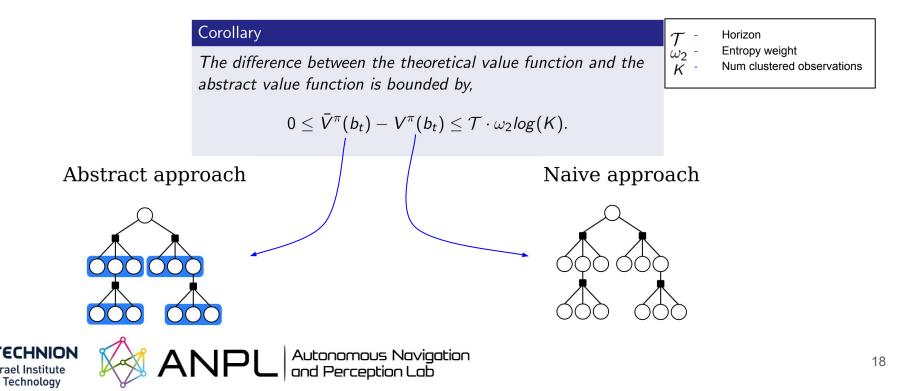
The new probability value is the aggregate average



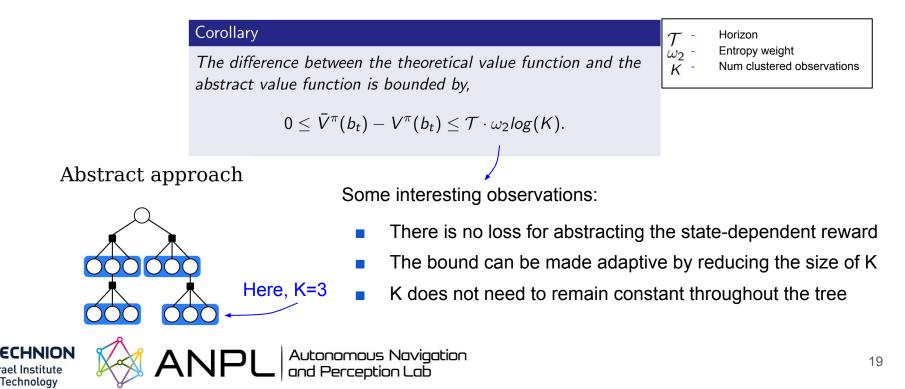




Using the abstract observation model, we derived analytical bounds compared to the non-abstract model,



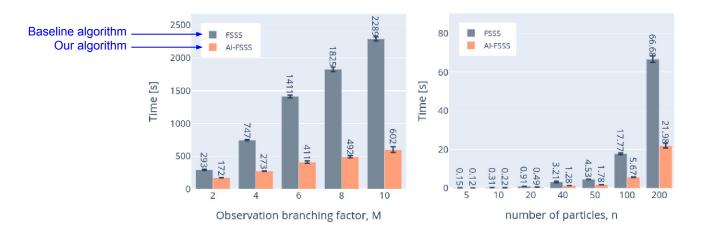
Using the abstract observation model, we derived analytical bounds compared to the non-abstract model,



Belief Based Rewards - Results

Speed-up for free

- Exact same solution
- A fraction of the planning time

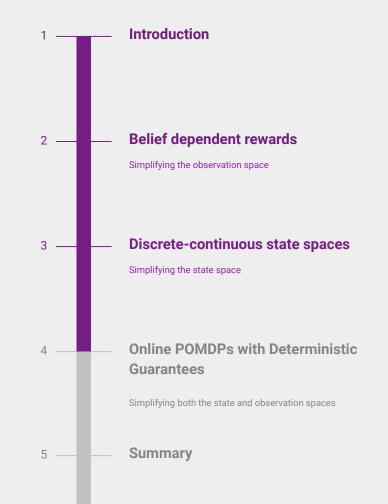




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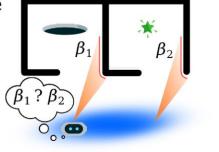
- Barenboim, M.; Shienman, M.; and Indelman, V. 2023., "Monte Carlo planning in hybrid belief POMDPs," IEEE Robotics and Automation Letters (RA-L).
- Barenboim, M.; Lev-Yehudi, I.; and Indelman, V. 2023. Data Association Aware POMDP Planning with Hypothesis Pruning Performance Guarantees. IEEE Robotics and Automation Letters (RA-L).

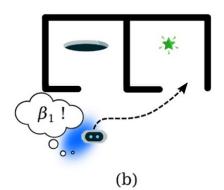


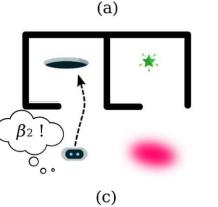
Continuous-Discrete State Spaces - Motivation

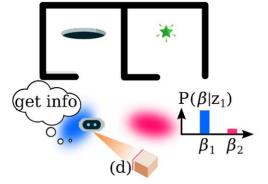
As an accompanying example, consider the case of ambiguous data associations:

- Uncertain observation source
- Optimality requires reasoning about different hypotheses









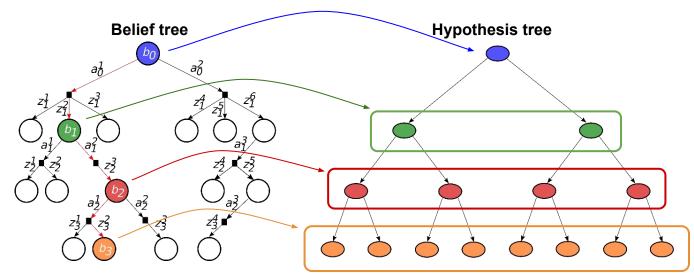




Continuous-Discrete State Spaces - The Challenge

Computing the reward function requires explicit knowledge of the hypotheses

However, the number of hypotheses may grow **exponentially** with the horizon!



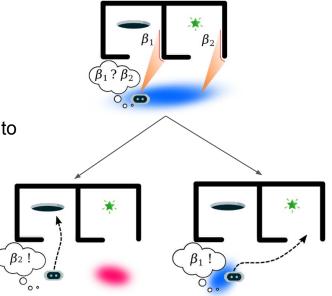


Continuous-Discrete State Spaces - Motivation

What others have done?

- Either implicitly assume known observation source
- Or prune hypotheses based on heuristics

It is not hard to show that a pruned set of hypotheses leads to a biased estimation







Continuous-Discrete State Spaces - Our Contributions (1)

Instead of computing all possible hypotheses, we utilize MCTS sampling and exploration approach. MCTS:

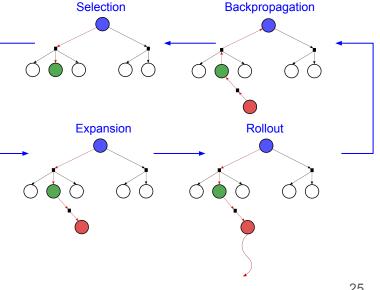
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- An MDP solver
- Uses UCT to tradeoff exploration-exploitation for actions

$$UCT(x_t, a_t) = \hat{Q}(x_t, a_t) + c \cdot \sqrt{\frac{\log N(x_t)}{n(x_t, a_t)}}$$

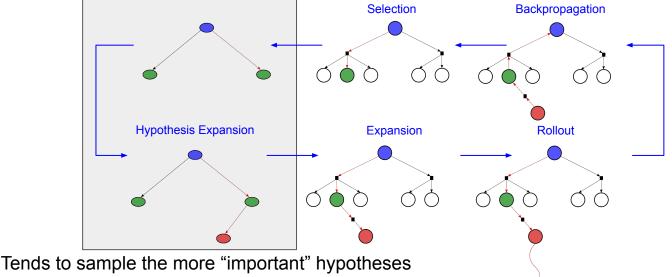
Given an action, samples the next state





Continuous-Discrete State Spaces - Our Contributions (1)

We add a layer that samples hypotheses via Monte-Carlo sampling



Can support belief-dependent rewards

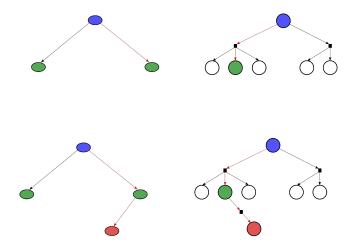


Continuous-Discrete State Spaces - Our Contributions (1)

We have derived a corresponding reward estimator, $\hat{\mathcal{R}}_X$, and have shown that it leads to an unbiased estimator,

Lemma

The sampled-based, state-dependent reward estimator, $\hat{\mathcal{R}}_X \triangleq \frac{1}{N} \sum_{i,j=1}^{N} \lambda_t^{i,j} \frac{1}{n_X} \sum_{k=1}^{n_X} r(X_t^{i,j,k}, a_t)$, is unbiased.

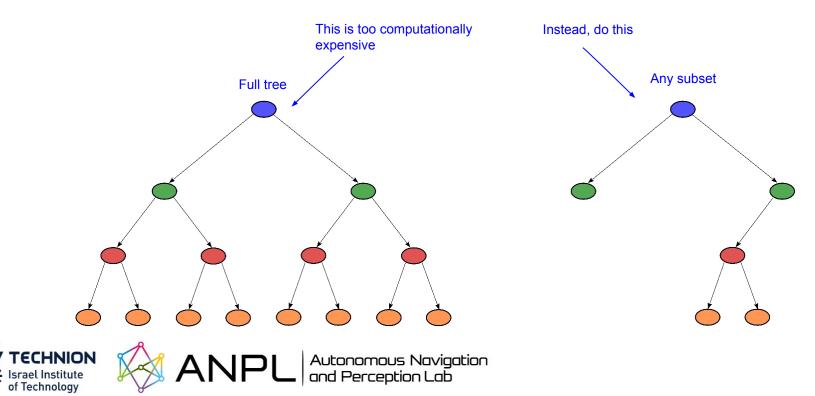






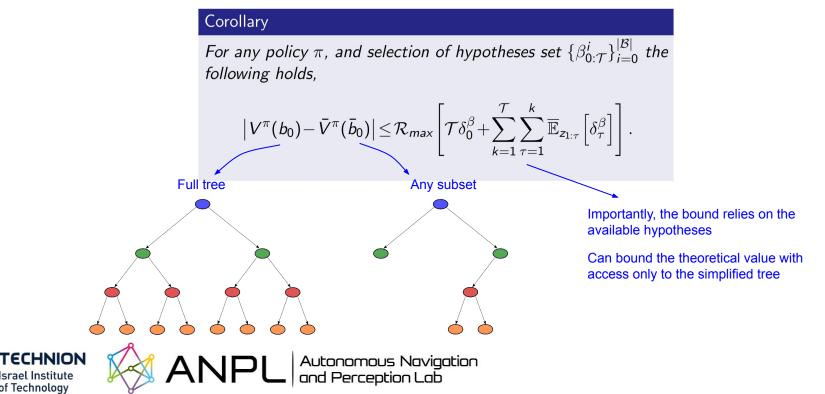
Continuous-Discrete State Spaces - Our Contributions (2)

Our second contribution bridges the gap between the full hypothesis tree and a simplified tree

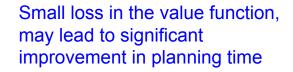


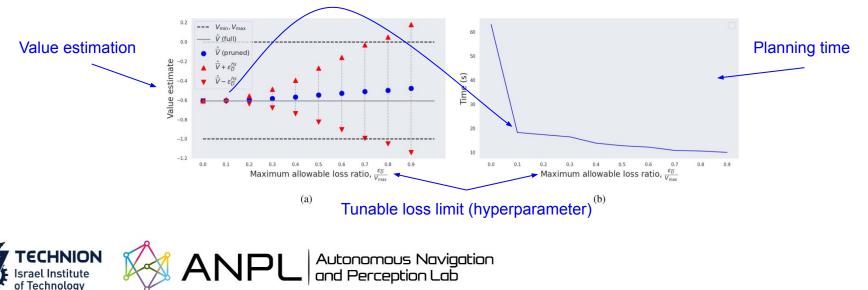
Continuous-Discrete State Spaces - Our Contributions (2)

Derived a deterministic bound to relate the full set of hypotheses to a subset thereof,



Continuous-Discrete State Spaces - Results





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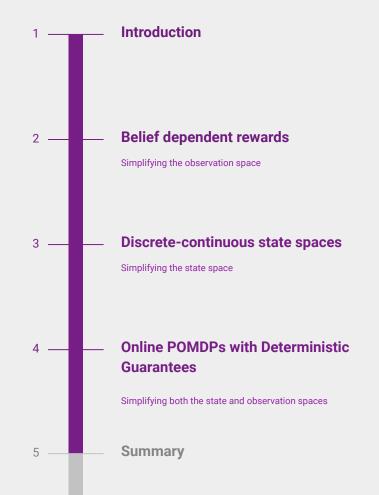
 M. Barenboim and V. Indelman. Online pomdp planning with anytime deterministic guarantees. In Advances in Neural Information Processing Systems, 2023

To be submitted:

Technoloav

 M. Barenboim and V. Indelman. Online pomdp planning with anytime deterministic guarantees - Extended version. To be submitted





POMDPs with Deterministic Guarantees - Motivation

Short reminder:

- POMDP is a formal framework for decision-making under uncertainty
- Finding an optimal policy is generally intractable
- Must resort to approximate solvers

Note - in this section we focus on discrete spaces

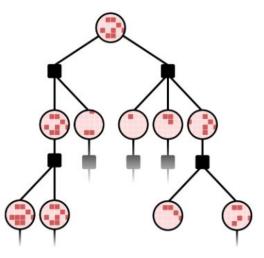




POMDPs with Deterministic Guarantees - Motivation

SOTA approximate solvers rely on sampling

They choose a subset of the state and observation spaces

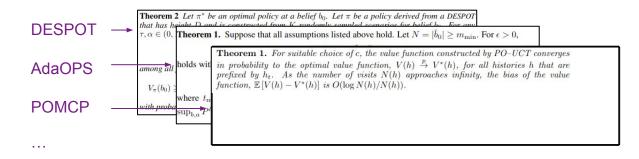






POMDPs with Deterministic Guarantees - The Challenge

Naturally, sampling comes with probabilistic theoretical guarantees



Can we get deterministic guarantees?



POMDPs with Deterministic Guarantees - Approach

In our work, we show that deterministic guarantees are indeed possible!

Given a POMDP:
$$\mathcal{M} = \langle \mathcal{X}, \mathcal{Z}, \mathcal{A}, \mathcal{P}_0, \mathcal{P}_T, \bar{\mathcal{P}}_Z, \mathcal{R}, \mathcal{T} \rangle$$

We define a simplified POMDP,

$$\begin{split} \bar{\mathcal{M}} &= \langle \bar{\mathcal{X}}, \bar{\mathcal{Z}}, \mathcal{A}, \bar{\mathcal{P}}_{0}, \bar{\mathcal{P}}_{\mathcal{T}}, \bar{\mathcal{P}}_{\mathcal{Z}}, \mathcal{R}, \mathcal{T} \rangle \\ \bar{\mathcal{X}}(H_{t}) &\subset \mathcal{X} \\ \bar{\mathcal{Z}}(H_{t}) &\subset \mathcal{Z} \\ \bar{\mathcal{I}}(H_{t}) &\subset \mathcal{Z} \\ \bar{\mathcal{I}}(H_{t}) &\subset \mathcal{Z} \\ \bar{\mathbb{P}}(x_{t+1} \mid x_{t}, a_{t}) \triangleq \begin{cases} \mathbb{P}(x_{t+1} \mid x_{t}, a_{t}) &, x_{t+1} \in \bar{\mathcal{X}}(H_{t+1}^{-}) \\ 0 &, otherwise \\ 0 &, otherwise \end{cases} \\ \bar{\mathbb{P}}(z_{t} \mid x_{t}) \triangleq \begin{cases} \mathbb{P}(z_{t} \mid x_{t}) &, z_{t} \in \bar{\mathcal{Z}}(H_{t}) \\ 0 &, otherwise \end{cases} \end{split}$$

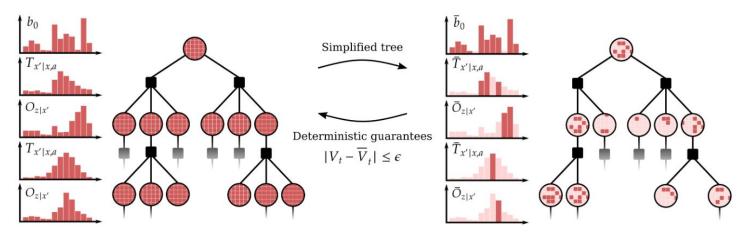


POMDPs with Deterministic Guarantees - Approach

With the simplified POMDP, we define a simplified value function,

$$ar{V}^{\pi}(ar{b}_t) riangleq r(ar{b}_t, \pi_t) + ar{\mathbb{E}}_{z_{t+1:\mathcal{T}}}\left[ar{V}^{\pi}(ar{b}_{t+1})
ight]$$

The formulation is flexible enough to allow any selection of the simplified state and observation spaces,





Derived upper and lower bounds for the optimal value function,

Lemma

Let \mathcal{A} be the set of actions and $\mathcal{U}_0^*(H_t)$, $\mathcal{L}_0^*(H_t)$ be the upper and lower bounds of node H_t . Then, the optimal value at the root is bounded by,

 $\mathcal{L}_0^{\star}(\mathcal{H}_0) \leq V^{\pi^*}(\mathcal{H}_0) \leq \mathcal{U}_0^{\star}(\mathcal{H}_0).$

- The bounds are easier to compute than the optimal value function
- The bounds shrink monotonically as the algorithm explores the tree
- Converge to the optimal value function
- This is the first work to our knowledge to provide deterministic guarantees for anytime online POMDPs



Importantly, the bounds can be calculated during planning.

How can we use them?

Pruning of sub-optimal branches

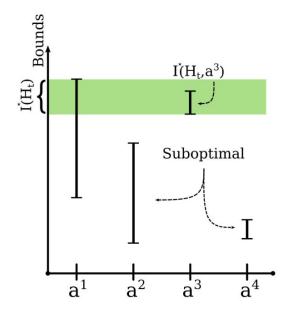
Made possible by the deterministic guarantees

Stopping criteria for the planning phase

Made possible by the deterministic guarantees

Finding the optimal solution in finite time

Without recovering the theoretical tree



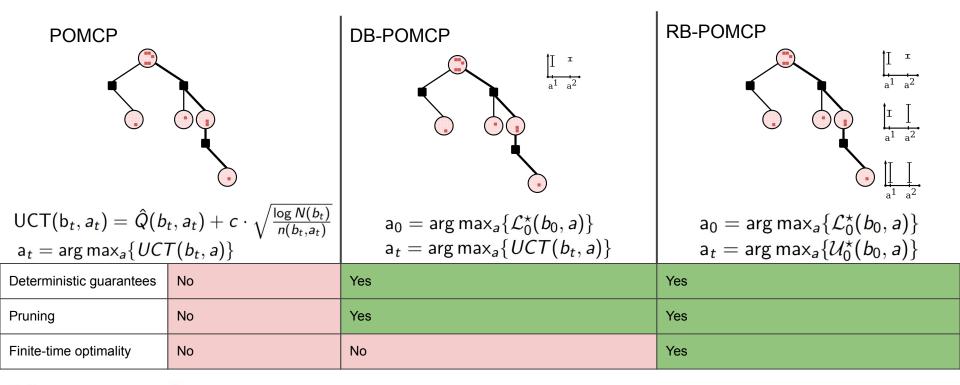




	Algorithm 1 ALGORITHM-A:	
Algorithm blueprint	function SEARCH	function SIMULATE $(h, d, \tau_d, \mathbb{P}_d)$
	1: while time permits do 2: Generate states x from b_0 .	1: if <i>d</i> = 0 then 2: return
	3: $\tau_0 \leftarrow x$	3: end if
	4: $\bar{\mathbb{P}}_0 \leftarrow b(x = \tau_0 \mid h_0)$	4: Select action a.
represents most SOTA algorithms	5: if $\tau_0 \notin \tau(h_0)$ then 6: $\overline{\mathbb{P}}(h_0) \leftarrow \overline{\mathbb{P}}(h_0) + \overline{\mathbb{P}}_0$	5: Generate next states and observations, x', z . 6: τ_d , $\mathbb{P}_{\tau} \leftarrow FWDUPDATE(ha, haz, \tau_d, \mathbb{P}_{\tau}, x')$
	7: end if	7: Select next observation z.
(cimilar atructura)	8: SIMULATE $(h_0, D, \tau_0, \overline{\mathbb{P}}_0)$.	8: SIMULATE $(haz, d-1, \tau_d, \bar{\mathbb{P}}_{\tau})$
(similar structure)	9: end while 1%: return	9: $BWDUPDATE(h, ha, d)$ 10: return
	10 return	
	function EWELLEDATE (he has a $\bar{\mathbb{D}}$ m')	
	function FWDUPDATE($ha, haz, \tau_d, \bar{\mathbb{P}}_{\tau}, x'$) 1: if $\tau_d \notin \tau(ha)$ then	function BWDUPDATE (h, ha, d) 1: $\epsilon(ha) = \gamma^{D-d} V_{max} s(\overline{\mathbb{P}}(h) - \overline{\mathbb{P}}(ha)) + \gamma^{D-d-1}$
	1: if $\tau_d \notin \tau(ha)$ then 2: $\tau(ha) \leftarrow \tau(ha) \cup \{\tau_d\}$	1: $\epsilon(ha) = \gamma^{D-d} V_{\max,d}(\overline{\mathbb{P}}(h) - \overline{\mathbb{P}}(ha)) + \gamma^{D-d-1}$
	1: if $\tau_d \notin \tau(ha)$ then 2: $\tau(ha) \leftarrow \tau(ha) \cup \{\tau_d\}$ 3: $\overline{R}(ha) \leftarrow \overline{R}(ha) + \overline{\mathbb{P}}_{\tau} \cdot r(x, a)$	1: $\epsilon(ha) = \gamma^{D-d} V_{\max,d}(\bar{\mathbb{P}}(h) - \bar{\mathbb{P}}(ha)) + \gamma^{D-d-1} V_{\max,d+1}(\bar{\mathbb{P}}(ha) - \sum_{z ha} \bar{\mathbb{P}}(haz))$
Cap attach our bounds to any such algorithm	1: if $\tau_d \notin \tau(ha)$ then 2: $\tau(ha) \leftarrow \tau(ha) \cup \{\tau_d\}$ 3: $\overline{R}(ha) \leftarrow \overline{R}(ha) + \overline{\mathbb{P}}_{\tau} \cdot r(x, a)$ 4: end if	1: $\epsilon(ha) = \gamma^{D-d} V_{\max,d}(\bar{\mathbb{P}}(h) - \bar{\mathbb{P}}(ha)) + \gamma^{D-d-1}$ $V_{\max,d+1}(\bar{\mathbb{P}}(ha) - \sum_{z ha} \bar{\mathbb{P}}(haz))$ 2: $U(ha) = \bar{R}(ha) + \gamma \sum_{z ha} U(haz) + \epsilon(ha)$
Can attach our bounds to any such algorithm	1: if $ au_{d} \notin au(ha)$ then 2: $ au(ha) \leftarrow au(ha) \cup \{ au_{d}\}$ 3: $\overline{R}(ha) \leftarrow \overline{R}(ha) + \overline{\mathbb{P}}_{\tau} \cdot r(x, a)$ 4: end if 5: $ au_{d} \leftarrow au_{d} \cup \{x'\}$ 6: $\overline{\mathbb{P}}_{\tau} \leftarrow \overline{\mathbb{P}}_{\tau} \cdot Z_{z x'} \cdot T_{x' x, a}$	1: $\epsilon(ha) = \gamma^{D-d} V_{\max,d}(\bar{\mathbb{P}}(ha) - \bar{\mathbb{P}}(ha)) + \gamma^{D-d-1}$ $V_{\max,d+1}(\bar{\mathbb{P}}(ha) - \sum_{z ha} \bar{\mathbb{P}}(haz))$ 2: $U(ha) = \bar{R}(ha) + \gamma \sum_{z ha} U(haz) + \epsilon(ha)$ 3: $L(ha) = \bar{R}(ha) + \gamma \sum_{z ha} L(haz) - \epsilon(ha)$
Can attach our bounds to any such algorithm	1: if $\tau_d \notin \tau(ha)$ then 2: $\tau(ha) \leftarrow \tau(ha) \cup \{\tau_d\}$ 3: $\overline{R}(ha) \leftarrow \overline{R}(ha) + \mathbb{P}_{\tau} \cdot r(x, a)$ 4: end if 5: $\tau_d \leftarrow \tau_d \cup \{x'\}$ 6: $\mathbb{P}_{\tau} \leftarrow \mathbb{P}_{\tau} \cdot Z_{z x'} \cdot T_{x' x,a}$ 7: if $\tau_d \notin \tau(haz)$ then	1: $\epsilon(ha) = \gamma^{D-d} V_{\max,d}(\bar{\mathbb{P}}(ha) - \bar{\mathbb{P}}(ha)) + \gamma^{D-d-1} V_{\max,d+1}(\bar{\mathbb{P}}(ha) - \sum_{ ha } \bar{\mathbb{P}}(haz))$ 2: $U(ha) = \bar{R}(ha) + \gamma \sum_{ ha } U(haz) + \epsilon(ha)$ 3: $L(ha) = \bar{R}(ha) + \gamma \sum_{ ha } L(haz) - \epsilon(ha)$ 4: $U(h) \leftarrow \max_{a'} \{U(ha')\}$
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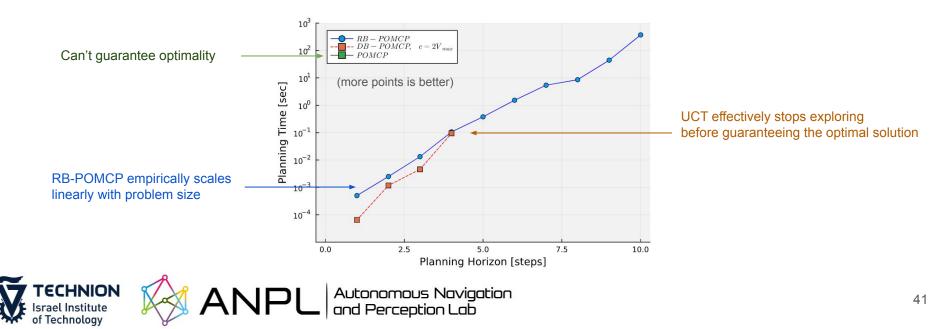




POMDPs with Deterministic Guarantees - Results

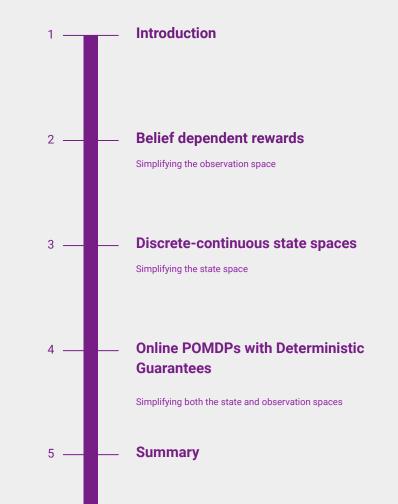
We compared the time it takes for the algorithms to find the optimal action

- SOTA algorithms excluded as they don't ensure optimal solutions
- Each point in the graph corresponds to the time it took to find the optimal action



Simplified POMDP Algorithms with Performance Guarantees

Summary







Thank you for listening!



